

Galois continued

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Recall: The Galois group of $K \subseteq F$ is

$$\text{Aut}_K F = \{ \varphi: F \rightarrow F \mid \varphi(k) = k \ \forall k \in K \}$$

We ended with:

Theorem: Suppose $K \subseteq F$ are fields, and $f \in K[x]$. Then if u is a root of f and $\sigma \in \text{Aut}_K F$, then $\sigma(u)$ is also a root of f .

As a consequence, suppose that u is a root of $f(x) \in K[x]$ and $F = K(u)$. Suppose that $\sigma \in \text{Aut}_K F$ and $\sigma(u) = v$. Then the lone equation $\sigma(u) = v$ completely determines $\sigma: F \rightarrow F$. This is because $\sigma: F \rightarrow F$ is a map of vector spaces, and a basis for F is $\{1, u, u^2, \dots, u^{n-1}\}$, and the action of σ on this basis is determined by $\sigma(u) = v$.

Since $\deg f = m$ implies f can have at most m distinct roots in F , this means

$$|\text{Aut}_K K(u)| \leq m$$

in fact $|\text{Aut}_K K(u)| \leq n$, where n is the number of distinct roots of f .

Example: If $F = K$, then $\text{Aut}_K F = \{\text{id}\}$. 90

However $\text{Aut}_K F = \{\text{id}\}$ does not imply $K = F$. For example, suppose u is the real cube root of 2.

Then $\text{Aut}_{\mathbb{Q}} \mathbb{Q}(u)$ contains automorphisms $\sigma: \mathbb{Q}(u) \rightarrow \mathbb{Q}(u)$ which permute the roots of $x^3 - 2$. But $\mathbb{Q}(u) \subseteq \mathbb{R}$, and the only roots of $x^3 - 2$ aside from u are complex. Thus ~~and~~ any $\sigma \in \text{Aut}_{\mathbb{Q}} \mathbb{Q}(u)$ must satisfy $\sigma(u) = u$ and thus be the identity (since a basis for $\mathbb{Q}(u)$ over \mathbb{Q} is $\{1, u, u^2\}$). So $\text{Aut}_{\mathbb{Q}} \mathbb{Q}(u) = \{\text{id}\}$.

Example: Let i denote a root of $x^2 + 1$, so that its two roots are $\pm i$. Set $\mathbb{C} = \mathbb{R}(i)$.

Then $\text{Aut}_{\mathbb{R}} \mathbb{C}$ has order at most two, since $x^2 + 1$ has two distinct roots.

There is a non-identity element of $\text{Aut}_{\mathbb{R}} \mathbb{C}$ given by $a + ib \mapsto a - ib$ (complex conjugation)

so $|\text{Aut}_{\mathbb{R}} \mathbb{C}| = 2$ and $\text{Aut}_{\mathbb{R}} \mathbb{C} \cong \mathbb{Z}_2$.

Example: Set $F = \mathbb{Q}(\sqrt{2}, \sqrt{3})$ and $K = \mathbb{Q}$. Then one can check that $\{1, \sqrt{2}, \sqrt{3}, \sqrt{6}\}$ is a basis for $\mathbb{Q}(\sqrt{2}, \sqrt{3})$ over \mathbb{Q} , so any $\sigma \in \text{Aut}_K F$ is determined by the two values $\sigma(\sqrt{2})$ and $\sigma(\sqrt{3})$.

But $\sigma(\sqrt{2}) = \pm \sqrt{2}$ (roots of $x^2 - 2$)

and $\sigma(\sqrt{3}) = \pm \sqrt{3}$ (roots of $x^2 - 3$).

Thus there are only 4 possibilities for σ . In fact, one can check that all possibilities are valid and that all give distinct automorphisms of order two.

Thus $|Aut_K F| = 4$, all elements of order 2

$$\Rightarrow Aut_K F = \mathbb{Z}_2 \times \mathbb{Z}_2.$$

Remarks • Given any finite group G , there exists an extension $K \subseteq F$ with $Aut_K F = G$.

• Open question: Fix a specific field K , say $K = \mathbb{Q}$. Which groups can arise from fixing the smaller field?

Idea of Galois theory: We will establish some sort of correspondence:

$$\left\{ \begin{array}{l} \text{fields } E \text{ with} \\ K \subseteq E \subseteq F \end{array} \right\} \overset{?}{\longleftrightarrow} \left\{ \begin{array}{l} \text{subgroups of} \\ Aut_K F \end{array} \right\},$$

eventually specializing to the case $[F : K] < \infty$. For now we keep it general, and have:

Theorem: Let F be an extension field of K , E an intermediate field and $H \leq Aut_K F$. Then

(a) $H' = \{h \in F \mid \sigma(h) = h \ \forall \sigma \in H\}$ is a subfield of F , and

(b) $E' = \{\sigma \in Aut_K F \mid \sigma(e) = e \ \forall e \in E\}$ is a subgroup of $Aut_K F$. (it is $Aut_E F$).

Proof: Easy exercise.

So we have a natural "first correspondence":

- Subgroups give fields by looking at the elements they fix
- Fields give subgroups by looking at the automorphisms that fix them.

Definition: The field H' in the Theorem above is the fixed field of $H \leq \text{Aut}_K F$.

We shall also use the "prime notation" on fields, i.e.

$\text{Aut}_E F \subset \text{Aut}_K F$ will be written $E' \subset \text{Aut}_K F$.

This notation then gives

$$\text{Aut}_K F = K' \leftarrow \{\text{id}\} \text{ and } F' = \{\text{id}\}.$$

On the other hand,

$$\{\text{id}\}' = \text{Aut}_K F \text{ with } (\text{Aut}_K F)' = ??$$

One might hope that $(\text{Aut}_K F)' = K$, so that everything aside from K is moved by at least one automorphism. But this is not the case: We already saw

$$\text{Aut}_{\mathbb{Q}} \mathbb{Q}(u) = \{\text{id}\} \text{ when } u \in \mathbb{R} \text{ is a root of } x^3 - 2,$$

$$\text{so } (\text{Aut}_{\mathbb{Q}} \mathbb{Q}(u))' = \mathbb{Q}(u) \text{ in that case.}$$

So in general, $(\text{Aut}_K F)'$ is some field E with $K \subseteq E \subseteq F$. $\text{Aut}_K F = \{\text{id}\}$ is the special case $E = F$, the other special case ($E = K$) gets a name:

Definition: Let $K \subseteq F$ be fields. If $(\text{Aut}_K F)' = K$ then F is said to be a Galois extension of K or is called Galois over K .

Example: We can easily check that \mathbb{C} is Galois over \mathbb{R} , since $\text{Aut}_{\mathbb{R}} \mathbb{C} = \mathbb{Z}_2$ and $a+ib \mapsto a-ib$ moves every element in $\mathbb{C} \setminus \mathbb{R}$. So $(\text{Aut}_{\mathbb{R}} \mathbb{C})' = \mathbb{R}$.

We now have all of the information and terminology needed to state the fundamental theorem of Galois theory:

Theorem: Suppose $K \subseteq F$ are fields and $[F:K] < \infty$.

If F is Galois over K then there is a correspondence

$$\left\{ \begin{array}{l} \text{fields } E \text{ with} \\ K \subseteq E \subseteq F \end{array} \right\} \iff \left\{ \begin{array}{l} \text{subgroups of} \\ \text{Aut}_K F \end{array} \right\}$$

given by $E \longmapsto E' = \text{Aut}_E F$
such that

(i) For any two intermediate fields L, M with $K \subseteq L \subseteq M \subseteq F$ we have

$$[M:L] = \frac{|M'|}{|L'|} = \frac{|M'|}{|L'|}$$

field dimension \curvearrowright \curvearrowleft group index

in particular $L \subseteq M \Rightarrow M' \subseteq L'$ and

$$|\text{Aut}_K F| = |\text{Aut}_K F : \{\text{id}\}| = [F : K].$$

(ii) F is Galois over every intermediate field E , but E is Galois over K if and only if E' is normal in $\text{Aut}_K F$. In this case

$$\text{Aut}_K F /_{E'} = \text{Aut}_K F /_{\text{Aut}_E F} \cong \text{Aut}_K E.$$

The fundamental theorem of Galois theory is "order-reversing": If

$$K \subset L \subset M \subset F \quad (\text{fields}) \quad \text{then}$$

$$\begin{array}{c} K' > L' > M' > F' \\ \text{"} & & \text{"} \\ \text{Aut}_K F & & \{\text{id}\} \end{array}$$

and vice versa: Given

$$\{\text{id}\} < H < J < \text{Aut}_K F, \quad \text{then}$$

$$\begin{array}{c} \{\text{id}\}' > H' > J' > (\text{Aut}_K F)' \\ \text{"} & & \text{"} \\ F & & K \end{array}$$

The "priming operations" outlined last class obey the following properties:

Lemma: Suppose $K \subseteq L \subseteq M \subseteq F$ are fields and $\{\text{id}\} \subset H, J \subset \text{Aut}_K F = G$ are groups. Then:

(i) $F' = \{\text{id}\}$ and $K' = G$

(ii) $\{\text{id}\}' = F$

(iii) $L \subseteq M \Rightarrow M' \subset L'$, and $H \subset J \Rightarrow J' \subseteq H'$

(iv) $L \subset L''$ and $H \subset H''$

(v) $L' = L'''$ and $H' = H'''$

Sketch: (i) - (iii) are definitions (or follow in one line from definitions)

To prove (iv), it is only a matter of unpacking the meaning of a double prime. To prove (v), note that $L \subset L'' \Rightarrow L''' \subset L'$, by applying (iii) and (iv) together. On the other hand, if we use the fact $H \subset H''$ with L' in place of H , then we get $L' \subset L'''$. Thus $L''' = L'$.

We argue similarly that $H' = H'''$.

Remark: • The containments $L \subset L''$ or $H \subset H''$ can easily be proper.

• Since F is Galois over K if $(\text{Aut}_K F)' = K$ and $K' = \text{Aut}_K F$, we get that F is Galois over K if and only if $(\text{Aut}_K F)' = K$, ie $K = K''$.

$$\text{"} \quad \text{"}$$

$$(K')' = K''$$

• By the same reasoning, if $K \subset E \subset F$ then F is Galois over E if and only if $E'' = E$.

Definition: If X is a field with $K \subset X \subset F$ or if X is a subgroup with $\{\text{id}\} \subset X \subset \text{Aut}_K F$, then X will be called closed if $X = X''$.

Theorem: If $K \subseteq F$ are fields, then there is a correspondence

$$\left\{ \begin{array}{l} \text{closed fields } E \\ K \subset E \subset F \end{array} \right\} \iff \left\{ \begin{array}{l} \text{closed subgroups } H \\ \{\text{id}\} < H < \text{Aut}_K F \end{array} \right\}$$

Given by $E \longmapsto E'$

which is one-to-one and onto.

Proof: It essentially follows from part (v) of the previous lemma:

Since $E' = E'''$ and $H' = H'''$, all primed objects are closed. Thus we only need to check that $H \longmapsto H'$ provides an inverse to $E \longmapsto E'$, at which point we can be certain that the correspondence is 1-1 and onto.

Begin our technical lemmas.

We will eventually show that in an algebraic Galois extension, all intermediate fields are closed.

When $[F:K] < \infty$, all subgroups of $\text{Aut}_K F$ are closed, too.

Lemma 2.8: Suppose $K \subset L \subset M \subset F$.

If $[M:L] < \infty$, then $|L':M'| \leq [M:L]$. In particular, if $[F:K] < \infty$ then

$$|K':F'| = |\text{Aut}_K F : \{\text{id}\}| = |\text{Aut}_K F| \leq [F:K].$$

Proof: First suppose $[M:L]=1$. Then $M=L$ and $L'=M'$ so $|L':M'| = [M:L]$, so we have a base case for an induction.

Now suppose the claim holds for $[M:L] < n$, and choose $u \in M \setminus L$. Since $[M:L] < \infty$, u is algebraic ^{consider $=n$.} over L with some minimal polynomial $f(x) \in L[x]$ with $\deg f = k > 1$. Then $[L(u):L] = k$ and since $[M:L] = n$, $[M:L(u)] = n/k$. I.e., we have:

$$\underbrace{L \subset L(u)}_k \subset \underbrace{L(u) \subset M}_{n/k}$$

$$\underbrace{\hspace{10em}}_n$$

and $L' \supseteq L(u)' \supseteq M'$, though we cannot yet say what the index of these subgroups in the larger group may be. To determine this, we consider cases:

Case 1: If $k < n$, then $1 < n/k < n$ and the induction hypothesis applies to both $[M:L(u)]$ and $[L(u):L]$ giving $[M:L(u)] \Rightarrow |L(u)':M'| \leq n/k$ and $|L':L(u)'| \leq k$ so

$$|L':M'| = |L':L(u)'| |L(u)':M'| \leq k \left(\frac{n}{k}\right) = n; \text{ done.}$$

Case 2: If $k = n$, then $[M:L(u)] = 1$ and $M = L(u)$.

This case requires work. Set

$$S = \{\tau M' \mid \tau \in L'\}, \text{ i.e. it's all left cosets of } M' \leq L'$$

and $T = \{\text{roots of } f(x) \text{ in } F\}$.

Construct an injective map $\psi: S \rightarrow T$ as follows:

(Completing this construction gives $|S| \leq |T|$, but

$|S| = |L' : M'|$ by definition and $|T| \leq n$ gives

$$|L' : M'| \leq n = [M : L], \text{ finishing the proof.}$$

Let $\tau M'$ be a left coset. Define ψ by

$$\psi(\tau M') = \tau(u).$$

First note that this map makes sense: Given $\tau \in L' = \text{Aut}_L F$, we know τ permutes the roots of $f(x)$, so $\tau(u)$ is indeed a root of f . Next, suppose $\tau M' = \tau' M'$. Then

$\tau = \tau' \sigma$ for some $\sigma \in M'$. Since $u \in M$, $\sigma(u) = u$ and therefore

$$\psi(\tau M') = \tau(u) = \tau' \sigma(u) = \tau'(u) = \psi(\tau' M'),$$

so ψ is well-defined.

Last, suppose $\psi(\tau M') = \psi(\tau' M')$. Then $\tau(u) = \tau'(u) \Rightarrow \tau^{-1} \tau'(u) = u$, so $\tau^{-1} \tau'$ fixes u . But then since we are in the situation $L(u) = M$, this means $\tau^{-1} \tau' : M \rightarrow M$ is the identity since it fixes the basis $\{1, u, \dots, u^{n-1}\}$. Thus $\tau^{-1} \tau' \in \text{Aut}_M F = M'$, so $\tau M' = \tau' M'$ and ψ is injective.

We have an analogous lemma relating priming to subgroups of the Galois group.

Lemma 2.9: Let F be an extension field of K and let H, J be subgroups of the group $\text{Aut}_K F$, with $H < J$. Then if $[J:H] < \infty$, we have $[H':J'] \leq [J:H]$.

Proof: Suppose that $[J:H] = n$ and $[H':J'] > n$. Then we can choose

- $u_1, u_2, \dots, u_{n+1} \in H'$ that are linearly independent over J'
- $\tau_1, \tau_2, \dots, \tau_n$ a set of coset representatives of H in J .

And consider the following system of equations:

$$\begin{array}{ccccccc} \tau_1(u_1)x_1 + \tau_1(u_2)x_2 + \dots + \tau_1(u_{n+1})x_{n+1} & = & 0 & & & & \\ \vdots & & \vdots & & \vdots & & \vdots \\ \tau_n(u_1)x_1 + \tau_n(u_2)x_2 + \dots + \tau_n(u_{n+1})x_{n+1} & = & 0 & & & & \end{array}$$

Because this system has $n+1$ variables and n equations, there is a nontrivial solution, in fact there are potentially many solutions. Choose a solution

$$x_1 = a_1, x_2 = a_2, \dots, x_{n+1} = a_{n+1}$$

with a minimal number of nonzero a_i . We can, up to reindexing variables and scaling our solution (recall any scalar multiple of a solution is also a solution)

assume that a_1, \dots, a_r are nonzero, a_{r+1}, \dots, a_{n+1} are zero, and $a_1 = 1$ (multiply by a_1^{-1} otherwise). Trick: we will find $\sigma \in J$ such that

$\sigma a_1, \sigma a_2, \sigma a_3, \dots, \sigma a_r, \sigma a_{r+1} = 0, \dots, \sigma a_{n+1} = 0$
is also a solution, with $\sigma a_2 \neq a_2$.

But then

$a_1 - \sigma a_1, a_2 - \sigma a_2, \dots, \text{etc}$

is a solution to the system of equations as well, and since $a_1 - \sigma a_1 = 1 - 1 = 0$ and $a_2 - \sigma a_2 \neq 0$ we have a nontrivial solution with fewer nonzero entries, contradiction to our minimal choice.

So to finish the proof we need only find σ .

Since τ_1, \dots, τ_n are coset reps of H in J , there's one τ_i that actually lies in H (a coset rep for $(\text{id}) \cdot H$). Then this τ_i , because it's in H , must fix every element in H' by definition. So $\tau_i(u_j) = u_j \forall j = 1, \dots, n+1$. Say $i=1$ so it's τ_1 fixing all u_j . Then the first equation in the system is

$$u_1 x_1 + u_2 x_2 + \dots + u_{n+1} x_{n+1} = 0.$$

Now since the a_i are a solution this means

$$\sum_{i=1}^{n+1} u_i a_i = 0$$

but the u_i are linearly independent over J' , so at least one of the a_i must not be in J' — let's say $a_2 \notin J'$.

Since $a_2 \notin J'$, by definition there's $\sigma \in J$ with $\sigma(a_2) \neq a_2$. So now consider

$$\begin{aligned} \sigma \tau_1(u_1)x_1 + \sigma \tau_1(u_2)x_2 + \dots + \sigma \tau_1(u_{n+1})x_{n+1} &= 0 \\ \vdots & \\ \sigma \tau_n(u_1)x_1 + \dots + \sigma \tau_n(u_{n+1})x_{n+1} &= 0, \end{aligned} \quad (*)$$

ie apply σ to the entire system we started with, and from this we also see that

$$\sigma a_1, \sigma a_2, \dots, \sigma a_{n+1}$$

is a solution to the new system above.

Claim: Although it looks different, aside from possibly re-ordering the equations, system (*) is the same as the original. Thus, we'll find σ is the automorphism we need.

(i) First, note that since $\{\tau_1, \dots, \tau_n\}$ are coset representatives of H in J and $\sigma \in J$, then $\{\sigma \tau_1, \dots, \sigma \tau_n\}$ are a complete set of coset representatives, too.

(If not, say $\sigma \tau_i H = \sigma \tau_j H$, so $\sigma \tau_i = \sigma \tau_j \gamma$ for some $\gamma \in H$.

Then $\tau_i = \tau_j \gamma \Rightarrow \tau_i H = \tau_j H$, contradiction).

(ii) Also note that if $\gamma, \beta \in \tau_j H$ for some j , then

$$\gamma = \tau_j h, \beta = \tau_j h' \text{ for } h, h' \in H \text{ and so } \forall i=1, \dots, n+1$$

$$\gamma(u_i) = \tau_j h(u_i) = \tau_j(u_i) = \tau_j h'(u_i) = \beta(u_i)$$

Since $u_i \in H'$