

MATH 2080, Lecture 5 § 0.4 continued.

Last day we defined equivalent sets, meaning there's a 1-1 onto (bijective) function between them. Today we consider a special case:

Example: Let $E \subset \mathbb{N}$ denote the set of even, positive integers. Define a map

$$f: E \rightarrow \mathbb{N}, f(2n) = n.$$

by dividing by two. Then it is easy to check that f is 1-1 and onto, since

$$f(2m) = f(2n) \Rightarrow m = n \quad (\text{so 1-1})$$

and given $n \in \mathbb{N}$, $f(2n) = n$. (so onto).

Thus $E \sim \mathbb{N}$.

Remark: This might seem strange - a set being equivalent to one of its proper subsets - but such is the nature of infinite sets.

Example: Define a function $f: \mathbb{N} \rightarrow \mathbb{Z}$ by

$$f(n) = \begin{cases} \frac{n}{2} & \text{if } n \text{ is } \del{\text{odd}} \text{ even} \\ \frac{1-n}{2} & \text{if } n \text{ is odd.} \end{cases}$$

Diagrammatically; f is going to have \mathbb{Z} as its image with

$$\{k \in \mathbb{Z} \mid k \leq 0\} = \text{image of odds}$$

$$\{k \in \mathbb{Z} \mid k > 0\} = \text{image of evens, and } f \text{ is 1-1 and onto.}$$

Let's prove this claim.

First, we check f is onto. Suppose $k \in \mathbb{Z}$ is positive, then $2k \in \mathbb{N}$ and $f(2k) = k$. On the other hand $\xrightarrow{\text{even}}$

If $k \leq 0$, then $1 - 2k \geq 1$ and $1 - 2k$ is odd, so we compute $f(1 - 2k) = \frac{1 - (1 - 2k)}{2} = k$.

So f is onto.

To show f is 1-1, suppose $f(n) = f(m)$.

If $f(n) = f(m) > 0$, then n and m must be even (because only evens map to positive numbers). Then we can calculate:

$$f(n) = f(m) \Rightarrow \frac{n}{2} = \frac{m}{2} \Rightarrow n = m.$$

On the other hand if $f(n) = f(m) \leq 0$, then n and m are both odd (since only odds map to numbers ≤ 0). Then

$$f(n) = f(m) \Rightarrow \frac{1-n}{2} = \frac{1-m}{2} \Rightarrow n = m$$

Therefore f is 1-1.

Thus $\mathbb{N} \sim \mathbb{Z}$.

Sets equivalent to \mathbb{N} get a special name.

Definitions: A set is countably infinite iff it is equivalent to \mathbb{N} . A set is finite if and only if it is empty or equivalent to $\{1, 2, \dots, n\}$ for some n . A set is countable if it is either finite or countably infinite. A set is infinite if it is not finite, and uncountable if it is not countable.

Question: which sets are countable? Is \mathbb{Q} countable? Is \mathbb{R} countable? We will answer these, eventually.

Theorem: Any subset of \mathbb{N} is countable.

Proof: If $S \subseteq \mathbb{N}$ is finite, it's countable. Suppose $S \subseteq \mathbb{N}$ is infinite. Then it's nonempty, by the well-ordering principle it has a smallest element, call it $f(1)$. Now after defining $f(1), f(2), \dots, f(k)$, define $f(k+1)$ to be the smallest element of $S \setminus \{f(1), f(2), \dots, f(k)\}$. This defines a function $f: \mathbb{N} \rightarrow S$. The function is 1-1 because of its construction. It is onto because for any $m \in S$, the set $\{x \in S \mid x < m\}$ is finite, say it has k elements. Then $\{f(1), f(2), \dots, f(k)\} = \{x \in S \mid x < m\}$ and $f(k+1) = m$. Therefore $S \sim \mathbb{N}$.

Corollary: Any subset of a countable set is countable.

Note that in order to prove countability of a set S , we no longer need to give a function $f: S \rightarrow \mathbb{N}$ that is 1-1 and onto, it's enough to give $f: S \rightarrow \mathbb{N}$ that's 1-1.

Theorem: If A and B are countable, so is $A \times B$.

Proof: If A and B are countable, there are 1-1 functions $f: A \rightarrow \mathbb{N}$ and $g: B \rightarrow \mathbb{N}$. Define $h: A \times B \rightarrow \mathbb{N}$ by $h(a, b) = \cancel{f(a)g(b)} 2^{f(a)} 3^{g(b)}$.

Then

$$h(a, b) = h(x, y)$$

$$\Rightarrow 2^{f(a)} 3^{g(b)} = 2^{f(x)} 3^{g(y)}$$

$\Rightarrow f(a) = f(x)$ and $g(b) = g(y)$ by unique factorization

$\Rightarrow a = x$ and $b = y$, so $(a, b) = (x, y)$ and

h is 1-1.

Example: \mathbb{Q} is countable.

Every element of \mathbb{Q} can be written $\frac{p}{q}$ where p, q have no common divisors and $q > 0$, moreover this representation is unique.

So define $f: \mathbb{Q} \rightarrow \mathbb{Z} \times \mathbb{N}$ by

$$f\left(\frac{p}{q}\right) = (p, q).$$

This function is 1-1 since the representation of each rational number is unique. The set $\mathbb{Z} \times \mathbb{N}$ is countable, therefore so is the image $f(\mathbb{Q})$, so \mathbb{Q} is countable.

We need one more method of making ^{new} countable sets from old, this will be used when we study continuity.

Theorem: Let $S \subseteq \mathbb{N}$ be a nonempty subset. Let $\{A_s\}_{s \in S}$ be a family of countable sets. Then $\bigcup_{s \in S} A_s$ is countable. ("A countable union of countable sets is countable").

Proof: For each $s \in S$, there's a 1-1 function

$f_s: A_s \rightarrow \mathbb{N}$. We need to build a 1-1 function

$f: \bigcup_{s \in S} A_s \rightarrow \mathbb{N} \times \mathbb{N}$ and we do it as follows:

Let $x \in \bigcup_{s \in S} A_s$, and let $m_x \in S$ be the smallest integer such that $x \in A_{m_x}$, this exists by the well-ordering principle. Define

$$f(x) = (m_x, f_{m_x}(x)).$$

We need to argue that f is 1-1. So suppose

$f(x) = f(y) = (m, n)$ for some $x, y \in \bigcup_{s \in S} A_s$.

Then x and y are both in A_m , because m is the smallest index such that $x \in A_m$ and $y \in A_m$.

Applying the 1-1 function $f_m: A_m \rightarrow \mathbb{N}$ to both x and y gives $f_m(x) = f_m(y) = n$, since f_m is 1-1 we get $x = y$. Thus f is 1-1. Therefore $\bigcup_{s \in S} A_s$ is countable.

Not all sets are countable. Two famous examples:

- $\mathcal{P}(\mathbb{N})$, the set of all subsets of \mathbb{N} is not countable (Exercise 37 proves this, to appear on Assl 2)
- \mathbb{R} is not countable, to be proved in class in mid October.

To study \mathbb{R} , it's actually enough to study $(0, 1)$, since:

Example: $\mathbb{R} \sim (0, 1)$.

Proof: Define $f: (0, 1) \rightarrow \mathbb{R}$ by $f(x) = \tan(\pi x - \frac{\pi}{2})$.

This is 1-1 and onto.

Challenge: Find a 1-1 onto function $f: (0, 1) \rightarrow \mathbb{R}$ without using trig.

Example of an uncountable set:

Let $X = \{\text{all functions } f: \mathbb{N} \rightarrow \mathbb{N}\}$. Suppose $\phi: \mathbb{N} \rightarrow X$ is any function, we'll show that ϕ cannot be surjective. Let $f_n: \mathbb{N} \rightarrow \mathbb{N}$ denote the function $\phi(n)$. Build a new function

$$g: \mathbb{N} \rightarrow \mathbb{N}$$

by defining $g(n) = f_n(n) + 1$ for all $n \in \mathbb{N}$.

Then $g \in X$, but g is not in the image of ϕ . It's not in the image of ϕ because:

Suppose it is. Then $g = f_k$ for some k . But when we plug k into both sides, we get

$$g(k) = f_k(k)$$

which is false since we set $g(k) = f_k(k) + 1$. So $g \notin \text{im } \phi$ and X is uncountable.

§0.5 Real numbers

It is possible to start with nothing but the natural numbers \mathbb{N} , and from them describe \mathbb{Z} , then \mathbb{Q} , and finally \mathbb{R} . This approach to studying the real numbers is extremely difficult (but interesting!) and is not how we shall approach the problem.

Instead, we will assume that there is a set denoted by \mathbb{R} , and functions

$$+ : \mathbb{R} \times \mathbb{R} \longrightarrow \mathbb{R} \quad \cdot : \mathbb{R} \times \mathbb{R} \longrightarrow \mathbb{R}$$

$$(x, y) \longmapsto x + y \quad (x, y) \longmapsto x \cdot y$$

and a relation $< \subset \mathbb{R} \times \mathbb{R}$ that together obey twelve rules: (called axioms).

1. $(x+y)+z = x+(y+z)$ and $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
2. $x+y = y+x$ and $x \cdot y = y \cdot x$
3. $x \cdot (y+z) = x \cdot y + x \cdot z$
4. There is exactly one element $0 \in \mathbb{R}$ such that $0+x = x$ for all $x \in \mathbb{R}$
5. For every $x \in \mathbb{R}$ there is exactly one $y \in \mathbb{R}$ such that $x+y=0$, we write $-x$ for such an element

6. There is exactly one element $1 \in \mathbb{R}$ such that $x \cdot 1 = x$ for all $x \in \mathbb{R}$.
7. For each $x \in \mathbb{R}$ with $x \neq 0$ there exists a unique $y \in \mathbb{R}$ such that $x \cdot y = 1$, we denote such a y by $1/x$. (or x^{-1})
8. $x < y$ implies $x + z < y + z$
9. $x < y$ and $y < z$ implies $x < z$
10. For all $(x, y) \in \mathbb{R} \times \mathbb{R}$ exactly one of $x < y$, $y < x$ or $x = y$ is true
11. $x < y$ and $z > 0 \Rightarrow xz < yz$.

These 11 properties make \mathbb{R} what is called an ordered field. But we want \mathbb{R} to be special in some way, so we add a 12th rule.

Definition: A set $S \subseteq \mathbb{R}$ is bounded from above (or below) if there is a number $M \in \mathbb{R}$ such that $x \leq M$ for all $x \in S$ (or $x \geq M$ for all $x \in S$).

The number M is called an upper bound (or lower bound) in this case. If S has an upper and lower bound, it's called a bounded set.

Example: Any finite set $S \subseteq \mathbb{R}$ is bounded, because it has a biggest and smallest element, which we denote by $\max S$ and $\min S$.

Example: The set $(1, \infty)$ is not bounded, but it is bounded below. All of the elements of $(-\infty, 1]$ are lower bounds. Obviously 1 is the "best" lower bound in some sense, though.

We introduce the idea of least upper bound and greatest lower bound to capture the idea of "best" upper and lower bounds.

Definition: Let $S \subseteq \mathbb{R}$. Then $a \in \mathbb{R}$ is a least upper bound for S if it is an upper bound for S and if b is also an upper bound for S , then $a \leq b$. Similarly $a \in \mathbb{R}$ is a greatest lower bound for S if it is a lower bound and if b is also a lower bound for S then $b \leq a$.

We abbreviate these terms as

$$a = \text{l.u.b. } S \text{ or } a = \sup S \text{ (supremum)}$$

and

$$a = \text{g.l.b. } S \text{ or } a = \inf S \text{ (infimum)}$$

respectively.

The 12th property of \mathbb{R} is then:

12. Every nonempty subset of \mathbb{R} that is bounded above has a least upper bound.

This 12th property distinguishes \mathbb{R} from other "ordered fields". For example, properties 1-11 are true in \mathbb{Q} , but 12 is not. Let's prove this.

Example: The set

$$X = \{x \in \mathbb{Q} \mid x^2 < 2\} \subset \mathbb{Q}$$

has no least upper bound in \mathbb{Q} .

Proof: The proof in full detail is very long (eg a full class at least) but the idea/sketch is this:

First note X is nonempty since $0 \in X$. Now suppose $\alpha \in \mathbb{Q}$ is a least upper bound for X . It is possible to show that we must have $\alpha^2 = 2$.

The idea of the proof is to prove that $\alpha^2 > 2$ forces α to be too large (it's not a least upper bound any more) and $\alpha^2 < 2$ forces α to not be an upper bound. So $\alpha^2 = 2$.

But now we ask: Is there $\alpha \in \mathbb{Q}$ with $\alpha^2 = 2$? No, here is the proof:

Suppose we choose $p, q > 0$ with no common divisors and $2 = \left(\frac{p}{q}\right)^2$.

Then $2q^2 = p^2$, so p^2 is even and therefore p is even (since p odd implies p^2 odd). So write $p = 2r$ for some integer r . Then

$$2q^2 = (2r)^2 = 4r^2$$

$$\Rightarrow q^2 = 2r^2,$$

so q is even, too. But we assumed p and q had no common divisors, a contradiction.

So property 12 distinguishes \mathbb{R} from \mathbb{Q} , indeed it distinguishes \mathbb{R} from all other ordered fields (making \mathbb{R} unique).

Now from these 12 properties we can prove all properties of \mathbb{R} that you're familiar with.

Theorem: Let $x, y, z \in \mathbb{R}$. Then:

- (i) If $x < y$, then $-y < -x$.
- (ii) $0 < 1$
- (iii) If $0 < x < y$ then $0 < \frac{1}{y} < \frac{1}{x}$
- (iv) If $x < y$ and $z < 0$ then $zx > zy$
- (v) $x^2 \geq 0$ (for all $x \in \mathbb{R}$).